

## ALGEBRA IN A GROTHENDIECK TOPOS: INJECTIVITY IN QUASI-EQUATIONAL CLASSES

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This paper gives a study of injectivity and some related notions in quasi-equational classes of algebras in an arbitrary Grothendieck topos  $\mathbf{E}$ . The main purpose is to describe the relationship between the class  $\text{mod } \Sigma$  of models of a set  $\Sigma$  of quasi-equations in the category of sets  $\text{Ens}$  and the corresponding class  $\text{mod}(\Sigma, \mathbf{E})$  of models of  $\Sigma$  in  $\mathbf{E}$  with respect to residual smallness, boundedness of essential extensions and injectivity. The basic nature of our results is that, for any given  $\Sigma$ , whatever holds in  $\text{Ens}$ , concerning these notions, also holds in  $\mathbf{E}$ . In particular, this substantially improves the earlier results of Howlett [5] regarding the existence of enough injectives in  $\text{mod}(\Sigma, \mathbf{E})$ .

### 1. Preliminaries

**1.1. Algebras in a category.** Let  $\mathbf{E}$  be a finitely complete category (in particular, it has a terminal object  $\mathbf{1}$ ). Given a family  $\tau = (n_\lambda)_{\lambda \in \Omega}$  of finite cardinal numbers  $n_\lambda$ , indexed by a set  $\Omega$ , recall that an algebra in  $\mathbf{E}$  is an entity  $A = (|A|, (\lambda_A)_{\lambda \in \Omega})$ , where  $|A|$  is an object of  $\mathbf{E}$ , called the underlying object of the algebra  $A$  and, for each  $\lambda \in \Omega$ , the  $\lambda$ -th operation  $\lambda_A: |A|^{n_\lambda} \rightarrow |A|$  of  $A$  is a morphism in  $\mathbf{E}$ ,  $n_\lambda$  being the arity of  $\lambda_A$ ; the family  $\tau = (n_\lambda)_{\lambda \in \Omega}$  is called the type of  $A$ . A homomorphism  $h: A \rightarrow B$  from an algebra  $A = (|A|, (\lambda_A)_{\lambda \in \Omega})$  to an algebra  $B = (|B|, (\lambda_B)_{\lambda \in \Omega})$  is a morphism  $|h|: |A| \rightarrow |B|$  such that the following diagram commutes, for each  $\lambda \in \Omega$ :

$$\begin{array}{ccc}
 |A|^{n_\lambda} & \xrightarrow{|h|^{n_\lambda}} & |B|^{n_\lambda} \\
 \lambda_A \downarrow & & \downarrow \lambda_B \\
 |A| & \xrightarrow{|h|} & |B|
 \end{array}$$

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The collection of all algebras (of the type  $\tau$ ) in  $\mathbf{E}$  and homomorphisms between them forms a category denoted by  $\text{Alg}(\tau)\mathbf{E}$  (or by  $\text{Alg}(\tau)$  if  $\mathbf{E} = \text{Ens}$ ).

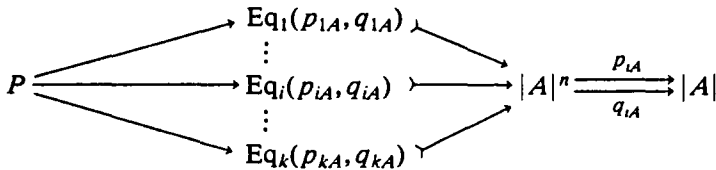
For  $A \in \text{Alg}(\tau)\mathbf{E}$  and any natural number  $n$ ,  $\mathbf{E}(|A|^n, |A|)$  can easily be made into an algebra of the type  $\tau$  in  $\text{Ens}$ , by defining the  $\lambda$ -th operation as

$$\lambda(\phi_1, \dots, \phi_{n_\lambda}) = \lambda_A \prod_{i=1}^{n_\lambda} \phi_i$$

for any  $\phi_i: |A|^n \rightarrow |A|$  ( $i = 1, \dots, n_\lambda$ ), where  $\prod_{i=1}^{n_\lambda} \phi_i$  is the morphism  $|A|^n \rightarrow |A|^{n_\lambda}$  determined by the  $\phi_i$ . Let  $F$  be the absolutely free algebra of the type  $\tau$  on a set  $X = \{x_1, \dots, x_n\}$  of  $n$  elements. Extend the map  $x_i \rightsquigarrow \text{pr}_i$  ( $\text{pr}_i: |A|^n \rightarrow |A|$  the projections) from  $X$  to  $\mathbf{E}(|A|^n, |A|)$  freely to a homomorphism  $\phi: F \rightarrow \mathbf{E}(|A|^n, |A|)$  and denote  $\phi(P)$  by  $P_A$  for any  $P \in F$ . For a quasi-equation

$$\sigma := \bigwedge_{i=1}^k (p_i = q_i) \rightarrow (p = q),$$

we say that  $A$  satisfies  $\sigma$ , written as  $A \models \sigma$ , iff the pullback



of the equalizers

$$\text{Eq}_i(p_{iA}, q_{iA}) \rightarrow |A|^n \begin{matrix} \xrightarrow{p_{iA}} \\ \xrightarrow{q_{iA}} \end{matrix} |A| \quad (i = 1, \dots, k)$$

factors through the equalizer

$$\text{Eq}(p_A, q_A) \rightarrow |A|^n \begin{matrix} \xrightarrow{p_A} \\ \xrightarrow{q_A} \end{matrix} |A|.$$

In particular,  $A$  satisfies an equation  $(p = q)$  iff  $p_A = q_A$ . The full subcategory of  $\text{Alg}(\tau)\mathbf{E}$  given by the class of all algebras in  $\mathbf{E}$  satisfying  $\Sigma$  will be denoted by  $\text{mod}(\Sigma, \mathbf{E})$  (or  $\text{mod } \Sigma$  if  $\mathbf{E} = \text{Ens}$ ) and is called a quasi-equational class (or an equational class if  $\Sigma$  is a set of equations). We note that, if  $\mathbf{E}$  is a topos, then the above definition of  $A \models \sigma$  coincides with the usual notion of satisfaction as defined for an arbitrary first order sentence in a topos. From now on  $\Sigma$  always denotes a set of quasi-equations.

Let  $k: \mathbf{E} \rightarrow \mathbf{F}$  be a functor, preserving finite limits; then  $k$  induces another functor  $f: \mathbf{A} \rightarrow \mathbf{B}$  in  $\text{Alg}(\tau)\mathbf{E}$ . Since  $k$  preserves finite limits, it preserves pullback and equalizer diagrams; and hence if  $\sigma$  is a quasi-equation and  $A \models \sigma$ , for  $A \in \text{Alg}(\tau)\mathbf{E}$ , then  $kA \models \sigma$ . We thus get a functor

$$\bar{k} | \text{mod}(\Sigma, \mathbf{E}) : \text{mod}(\Sigma, \mathbf{E}) \rightarrow \text{mod}(\Sigma, \mathbf{F})$$

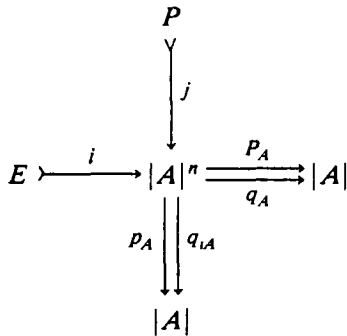
for any given set  $\Sigma$  of quasi-equations.

**1.2. Lemma.** *Let  $\mathbf{E}$  have a set  $\Phi$  of generators. Then, for any  $A \in \text{Alg}(\tau)\mathbf{E}$  and any set  $\Sigma$  of quasi-equations,  $A \in \text{mod}(\Sigma, \mathbf{E})$  iff  $\bar{h}_G(A) \in \text{mod} \Sigma$  for each  $G \in \Phi$ , where  $h_G = \mathbf{E}(G, -)$ .*

**Proof.** That  $A \in \text{mod}(\Sigma, \mathbf{E})$  implies  $\bar{h}_G(A) \in \text{mod} \Sigma$  is clear, by what has been discussed above. Conversely, let

$$\sigma := \bigwedge_{i=1}^k (p_i = q_i) \rightarrow (p = q)$$

be any quasi-equation and  $A \in \text{Alg}(\tau)\mathbf{E}$  such that  $\bar{h}_G(A) \models \sigma$  for all  $G \in \Phi$ . Consider the diagram



with  $P$  the pullback of the equalizers of the pairs  $(p_{iA}, q_{iA})$ . By the hypothesis on  $A$ ,  $h_G(p_A j) = h_G(q_A j)$  for all  $G \in \Phi$  which this implies that  $p_A j = q_A j$ , the latter because the set  $\Phi$  of generators is collectively faithful, and hence  $j$  factors through  $i$ , by definition of equalizers; thus  $A \models \sigma$ .

**1.3. Algebras in a Grothendieck topos.** Let  $\mathbf{C}$  be a small category. The category  $\hat{\mathbf{C}} = \text{Ens}^{\mathbf{C}^*}$  ( $\mathbf{C}^*$  the dual category of  $\mathbf{C}$ ) is called the category of presheaves on  $\mathbf{C}$ . Recall that, up to equivalence, a category  $\mathbf{E}$  is a *Grothendieck topos* iff it is a full subcategory of some  $\hat{\mathbf{C}}$  for which there is a reflection functor  $R: \hat{\mathbf{C}} \rightarrow \mathbf{E}$  (i.e. a left adjoint to the inclusion function  $\mathbf{E} \xrightarrow{i} \hat{\mathbf{C}}$ ) which is also left exact (i.e. preserves finite limits) (Johnstone [6], p. 12 and p. 105). Let

$$\begin{array}{ccc}
 & i & \\
 \mathbf{E} & \xrightleftharpoons[R]{} & \hat{\mathbf{C}}
 \end{array}$$

be a Grothendieck topos. The discussion in the last sections, in particular, shows that the category  $\text{Alg}(\tau)\hat{\mathbf{C}}$  is isomorphic to the category of all  $\text{Alg}(\tau)$ -valued presheaves on  $\mathbf{C}$ ; and since  $R$  preserves finite limits, it can be lifted to  $\bar{R}: \text{mod}(\Sigma, \hat{\mathbf{C}}) \rightarrow \text{mod}(\Sigma, \mathbf{E})$  (we denote  $\bar{R}$  by the same letter  $R$ ). Moreover, since  $\mathbf{E}$  has a set of generators, namely  $\{R(h_U): U \in \mathbf{C}\}$ , where  $h_U = (-, U)$  is the presheaf represented by  $U$ , Lemma 1.3 implies that  $A \in \text{mod}(\Sigma, \mathbf{E})$  iff  $AU \in \text{mod} \Sigma$  for all  $U \in \mathbf{C}$ .

A natural question to ask would be, what is the relationship between the

behaviour of a certain classical algebraic notion in  $\text{mod } \Sigma$  and in  $\text{mod}(\Sigma, \mathbf{E})$ . In this paper, we primarily consider the notion of injectivity, which has been extensively investigated for the case of equational classes of algebras in  $\text{Ens}$  (e.g. [1], [7]), and show that the properties of  $\text{mod } \Sigma$ , regarding this notion, survive the passage to  $\text{mod}(\Sigma, \mathbf{E})$ , for a set  $\Sigma$  of quasi-equations and an arbitrary Grothendieck topos  $\mathbf{E}$ . For example, we prove that  $\text{mod}(\Sigma, \mathbf{E})$  has enough injectives iff it is residually small and pushouts transfer monomorphisms, which is a counterpart of a result for equational classes of algebras in  $\text{Ens}$ ; further, we show that  $\text{mod}(\Sigma, \mathbf{E})$  has enough injectives iff  $\text{mod } \Sigma$  has enough injectives.

2. The adjointness of  $\hat{\mathbf{C}}$

2.1. Here, we construct a pair of adjoint functors

$$\hat{\mathbf{C}} \begin{matrix} \xrightarrow{G} \\ \xleftarrow{H} \end{matrix} \text{Ens}^{|\mathbf{C}|}$$

with  $G$  a left adjoint to  $H$ , where  $|\mathbf{C}|$  denotes the set of object of  $\mathbf{C}$ .

Define  $G$  by  $GP = (PU)_{U \in \mathbf{C}}$ , for  $P \in \hat{\mathbf{C}}$ , and for any map  $f: P \rightarrow Q$  in  $\hat{\mathbf{C}}$ ,  $Gf = (f_U)_{U \in \mathbf{C}}$ . That  $G$  is a functor is easily checked.

Define  $H$  by: for any  $B = (B_V)_{V \in \mathbf{C}}$  in  $\text{Ens}^{|\mathbf{C}|}$ ,  $HB = \bar{B}$  is the presheaf given as follows:

$$BU = \prod_{V \in \mathbf{C}} B_V^{(V, U)}$$

for each  $V \in \mathbf{C}$  and to define  $\bar{B}s: \bar{B}W \rightarrow \bar{B}U$ , for each  $s: U \rightarrow W$ , notice that  $s$  induces a natural transformation  $s^*: h_U \rightarrow h_W$  with components  $s_V^*$  given by  $s_V^*(t) = st$  ( $V \in \mathbf{C}$ ), for  $t: V \rightarrow U$ ; then  $s^*$  induces a map  $s_V^B: B_V^{(V, W)} \rightarrow B_V^{(V, U)}$  (composition from the left). Hence, define  $\bar{B}s = \prod_{V \in \mathbf{C}} s_V^B$ . One easily checks that  $\bar{B}$  is a presheaf. To complete the definition of the functor  $H$ , let  $f = (f_V)_{V \in \mathbf{C}}$  be any map from  $B = (B_V)_{V \in \mathbf{C}}$  to  $C = (C_V)_{V \in \mathbf{C}}$ . For any pair  $V$  and  $U$  in  $\mathbf{C}$ , define  $f_V^U: B_V^{(V, U)} \rightarrow C_V^{(V, U)}$  by  $f_V^U(\alpha) = f_V \alpha$  for any  $\alpha: (V, U) \rightarrow B_V$ . Define  $\bar{f} = Hf$  by  $\bar{f}_U = \prod_V f_V^U$ , ( $U \in \mathbf{C}$ ). That  $\bar{f}$  is a natural transformation and  $H$  is actually a functor is easily checked.

Next, we define two natural transformations  $\eta: 1 \rightarrow HG$  and  $\varepsilon: GH \rightarrow 1$ . To define  $\eta$ , let  $P \in \hat{\mathbf{C}}$ ; combining the maps  $PU \rightarrow PV^{(V, U)}$  given by  $a \rightsquigarrow a_V$  with  $a_V(s) = (Ps)a$ , for  $a \in PU$  and  $s: V \rightarrow U$ , we get a map

$$\eta_{PU}: PU \rightarrow \prod_{V \in \mathbf{C}} PV^{(V, U)} = (HGP) \quad (U \in \mathbf{C}).$$

Define  $\eta_P: P \rightarrow HGP$  by  $(\eta_P)_U = \eta_{PU}$ . It is not hard to check that each  $\eta_P$  and then  $\eta$  are natural transformations. Note that, for  $V = U$  and  $s = 1_U$ ,  $\bar{a}(s) = a$  and hence  $\eta_P$  is a monomorphism for each  $P \in \hat{\mathbf{C}}$ . We define  $\varepsilon$  such that, for each  $B = (B_V)_{V \in \mathbf{C}}$ ,  $\varepsilon_B$  has the composite

$$e_{pr_u}: \prod_V B_V^{(V, U)} \rightarrow B_U^{(U, U)} \rightarrow B_U$$

as its  $U$ -th component ( $U \in \mathbf{C}$ ) and  $e(\alpha) = \alpha(1_U)$  for  $\alpha : (U, U) \rightarrow B_U$ .

**2.2. Proposition.**  $G$  is a left adjoint to  $H$ , with  $\eta$  and  $\varepsilon$  as front and back adjunctions.

**Proof.** It only remains to show that

$$1_G = G \xrightarrow{G \circ \eta} GHG \xrightarrow{\varepsilon \circ G} G$$

and

$$1_H = H \xrightarrow{\eta \circ H} HGH \xrightarrow{H \circ \varepsilon} H,$$

that is to show, for any  $A \in \hat{\mathbf{C}}$  and  $B = (B_V)_{V \in \mathbf{C}}$ ,  $\varepsilon_{GA} G \eta_A = 1_{GA}$  and  $H \varepsilon_B \eta_{HB} = 1_{HB}$ . Let  $U \in \mathbf{C}$  and  $a \in AU$ , we have  $(GA)U \rightarrow (GHGA)U \rightarrow (GA)U$  given by  $a \rightsquigarrow (\tilde{a}_V)_{V \in \mathbf{C}} \rightsquigarrow e(\tilde{a}_V)$ , but  $e(\tilde{a}_V) = \tilde{a}_V(1_U) = A(1_U)(a) = a$ , and hence  $\varepsilon_{GA} G \eta_A = 1_{GA}$ . Similarly, one can show  $H \varepsilon_B \eta_{HB} = 1_{HB}$ .

**2.3. Remark.** By the definition of  $G$  and since  $H$  is a right adjoint, they both preserve finite limits, and hence can be lifted to

$$\text{mod}(\Sigma, \hat{\mathbf{C}}) \xrightleftharpoons[H]{G} \text{mod } \Sigma^{|\mathbf{C}|}$$

with  $\hat{G}$  a left adjoint to  $\hat{H}$ .

**2.4. Remark.** The particular case of the above construction for a monoid  $\mathbf{M}$  as the category  $\mathbf{C}$ , and hence the category of  $\mathbf{M}$ -sets as  $\hat{\mathbf{C}}$ , is due to Berthiaume [4].

### 3. Residual smallness

**3.1. Definition.** A monomorphism  $h : A \rightarrow B$  in a category  $\mathbf{K}$  is called *essential* iff, for any  $g : B \rightarrow C$  in  $\mathbf{K}$ , whenever  $gh$  is a monomorphism, then so is  $g$ .

Let  $\mathbf{E}$  be a Grothendieck topos and  $\Sigma$  a set of quasi-equations. Note that, as in the case of  $\text{Ens}$ ,  $\text{mod}(\Sigma, \mathbf{E})$  is closed under direct limits in  $\text{Alg}(\tau)\mathbf{E}$  because any colimit in  $\mathbf{E}$  is formed by first forming it in  $\hat{\mathbf{C}}$  and then reflecting it to  $\mathbf{E}$  where the latter preserves the conditions for  $A \models \Sigma$ .

**3.2. Lemma.** In  $\text{mod}(\Sigma, \mathbf{E})$ ,

- (i) any composite of essential monomorphisms is an essential monomorphism, and
- (ii) any direct limit of essential monomorphisms is an essential monomorphism.

**Proof.** (i) is trivial. To prove (ii), let  $f : A \rightarrow \varinjlim B_\alpha$  be a direct limit in  $\text{mod}(\Sigma, \mathbf{E})$  of

essential monomorphisms  $f_\alpha: A \rightarrow B_\alpha$  and diagram maps  $g_{\alpha\beta}: B_\alpha \rightarrow B_\beta$  ( $\beta \geq \alpha$ ). Since each  $f_\alpha$  is a monomorphism,  $f$  is a monomorphism, because this is true in  $\mathbf{Ens}$  and hence in  $\mathbf{\hat{C}}$ , and the reflection to  $\mathbf{E}$  preserves monomorphisms by exactness. To show that  $f$  is essential, let  $hf: A \rightarrow D$  be a monomorphism, with  $D$  in  $\text{mod}(\Sigma, \mathbf{E})$ . Then essentialness of  $f_\alpha$  implies that all  $hg_\alpha$  are monomorphisms, and hence  $h$  is a monomorphism; the latter again because of the way direct limits are formed in  $\text{mod}(\Sigma, \mathbf{E})$ . This proves that  $f$  is an essential monomorphism.

**3.3.** The following is a counterpart of (E3) in [1].

**Lemma.** *In  $\text{mod}(\Sigma, \mathbf{E})$ , for any monomorphism  $h: A \rightarrow B$  there exists a homomorphism  $g: B \rightarrow C$  with  $gh$  an essential monomorphism.*

**Proof.** Take all the congruences  $\Theta$  on  $B$  such that  $B/\Theta \in \text{mod}(\Sigma, \mathbf{E})$  and  $A \xrightarrow{vh} B/\Theta$  ( $v: B \rightarrow B/\Theta$  the quotient map) is a monomorphism. Then by the observation in 3.1 and the exactness argument in the proof of 3.2, any join of a chain of such congruences is again such a congruence; and hence, there exists a maximal such congruence, say  $\Theta_0$ . Maximality of  $\Theta_0$  then implies that  $A \rightarrow B/\Theta_0$  is essential.

**3.4. Corollary.** *In  $\text{mod}(\Sigma, \mathbf{E})$ , an algebra  $A$  is an absolute retract iff it has no proper essential extension.*

**Proof.** ( $\Rightarrow$ ) If  $f: A \rightarrow B$  is an essential monomorphism and  $h: B \rightarrow A$  is a retraction, then, by essentialness of  $f$ ,  $h$  is a monomorphism, and hence  $A \cong B$ .

( $\Leftarrow$ ) Given any monomorphism  $f: A \rightarrow B$ , continue it to an essential monomorphism  $A \xrightarrow{gf} C$ , by the last lemma. By hypothesis on  $A$ ,  $gf$  is an isomorphism and then  $(gf)^{-1}g$  is the desired retraction.

**3.5. Definition.** A category  $\mathbf{K}$  is called *residually small* iff, it has a set of cogenerators.

**3.6. Definition.** A category  $\mathbf{K}$  is called *essentially bounded* iff, for each  $A \in \mathbf{K}$  there exists, up to isomorphism, only a set of essential extensions in  $\mathbf{K}$ .

**3.7. Lemma.** *For any well powered category with products and a set  $\Phi$  of generators, residual smallness implies essential boundedness.*

**Proof.** Let  $h: A \rightarrow B$  be any essential monomorphism, and embed  $B \xrightarrow{e} \prod_{\alpha \in I} C_\alpha$ , for  $C_\alpha$  from a suitable set of cogenerators. Then, for any  $G \in \Phi$  and a pair of distinct maps  $G \xrightarrow{s} A$ , we have  $ehs \neq eht$ , and hence  $p_\alpha ehs \neq p_\alpha eht$ , for some projection  $p_\alpha: \prod C_\alpha \rightarrow C_\alpha$ . Pick  $\alpha_{st}$  as one such, then  $A \rightarrow \prod_{\beta \in J} C_\beta$  is a monomorphism, where  $J = \{\alpha_{st}: s \neq t: G \rightarrow A\}$  and  $\text{Card } J \leq \text{Card } \bigcup_{G \in \Phi} (G, A)^2$ . Essentialness of  $h$  implies that  $B \rightarrow \prod_J C_\beta$  is a monomorphism, and since there exists only a set of products  $\prod_J C_\beta$ , we are done.

**3.8. Lemma.** *For  $\text{mod}(\Sigma, \mathbf{E})$ , essential boundedness implies residual smallness.*

**Proof.** For any  $A \in \text{mod}(\Sigma, \mathbf{E})$ , take all  $B_\alpha \leq A$ , generated by

$$\alpha : U \amalg U \rightarrow |A| \quad (U \in \mathbf{C})$$

and then continue them to essential extensions

$$B_\alpha \xrightarrow{i_\alpha} A \xrightarrow{f_\alpha} C_\alpha,$$

by Lemma 3.3. The homomorphism  $\amalg f_\alpha : A \rightarrow \amalg C_\alpha$  is a monomorphism, for otherwise there exists some  $\alpha$  with  $B_\alpha \rightarrow A \rightarrow \amalg C_\alpha$  not a monomorphism which is a contradiction to the fact that all  $B_\alpha \rightarrow C_\alpha$  are monomorphisms. Since there exists, up to isomorphisms, only a set of  $B$  generated by some  $U \amalg U \rightarrow |B|$ , and only a set of essential extensions of those  $B$ , any set representing, up to isomorphism, all essential extensions of such  $B$  is a cogenerating set; hence, we are done.

**3.9. Corollary.** *For  $\text{mod}(\Sigma, \mathbf{E})$ , essential boundedness is equivalent to residual smallness.*

**Proof.** One way this is true by the last lemma, and since  $\text{mod}(\Sigma, \mathbf{E})$  has a set of generators, namely the  $\text{mod}(\Sigma, \mathbf{E})$ -free algebras on the reflection of representable presheaves  $h_U$  ( $U \in \mathbf{C}$ ), Lemma 3.7 implies the converse.

**3.10. Proposition.**  *$\text{mod}(\Sigma, \mathbf{E})$  is residually small iff  $\text{mod } \Sigma$  is residually small.*

**Proof.** ( $\Rightarrow$ ) Consider the following pair of adjoint functors:

$$\text{mod}(\Sigma, \mathbf{E}) \begin{matrix} \xrightarrow{\Gamma} \\ \xleftarrow{\Delta} \end{matrix} \text{mod } \Sigma$$

where  $\Gamma = (\mathbf{1}, -)$  and  $\Delta$  left exact, left adjoint to  $\Gamma$ ; in fact  $\Delta$  is the composite

$$\text{mod } \Sigma \xrightarrow{\Delta_0} \text{mod } \Sigma^{|C|} \xrightarrow{|H|} \text{mod}(\Sigma, \hat{\mathbf{C}}) \xrightarrow{R} \text{mod}(\Sigma, \mathbf{E})$$

where  $\Delta_0$  takes  $\text{mod } \Sigma$  to the constant families induced by  $|C|$  and  $\hat{H}$  as in Section 2. One can then easily check that the functor  $\Gamma$  transfers the set of cogenerators, and hence  $\text{mod } \Sigma$  is residually small.

( $\Leftarrow$ ) Let  $\text{mod } \Sigma$  be residually small. Then so is any  $\text{mod } \Sigma' = \text{mod}(\Sigma, \text{Ens}^I)$ . The functor  $\hat{H}$  transfers the set of cogenerators of  $\text{mod}(\Sigma, \text{Ens}^{|C|})$  to a set of cogenerators in  $\text{mod}(\Sigma, \hat{\mathbf{C}})$ , and hence the latter is essentially bounded, by Lemma 3.7. Since monomorphisms in  $\text{mod}(\Sigma, \mathbf{E})$  are also monomorphisms in  $\text{mod}(\Sigma, \hat{\mathbf{C}})$  and the reflection functor preserves monomorphisms, essential monomorphisms in  $\text{mod}(\Sigma, \mathbf{E})$  are also essential in  $\text{mod}(\Sigma, \hat{\mathbf{C}})$ , and hence  $\text{mod}(\Sigma, \mathbf{E})$  is also essentially bounded. Then  $\text{mod}(\Sigma, \mathbf{E})$  is residually small, by Lemma 3.8.

4. Injective algebras in  $\text{mod}(\Sigma, \mathbf{E})$

4.1. Definition. In any category  $\mathbf{K}$ , pushouts transfer monomorphisms iff, for any pushout diagram

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ \downarrow u & & \downarrow v \\ C & \xrightarrow{g} & D \end{array}$$

whenever  $f$  is a monomorphism, then  $g$  is also a monomorphism. If  $\mathbf{K}$  has pushouts, one can easily check that this is equivalent to saying that, any diagram

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ \downarrow u & & \\ C & & \end{array}$$

with  $f$  a monomorphism can be completed to a commutative diagram

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ \downarrow u & & \downarrow v \\ C & \xrightarrow{g} & D \end{array}$$

with  $g$  a monomorphism. This is the condition (E4) in [1].

4.2. Proposition. Pushouts transfer monomorphisms in  $\text{mod}(\Sigma, \mathbf{E})$  iff they do in  $\text{mod } \Sigma$ .

Proof. ( $\Rightarrow$ ) Using the pair of adjoint functors  $\Delta \dashv \Gamma$  given in the proof of Proposition 3.10 and the fact that  $\Delta$  is faithful, one can easily check this.

( $\Leftarrow$ ) If the diagram

$$\begin{array}{ccc} A & \xrightarrow{f} & B \\ \downarrow u & & \downarrow v \\ C & \xrightarrow{g} & D \end{array}$$

is a pushout in  $\text{mod}(\Sigma, \mathbf{E})$  with  $f$  a monomorphism, then, by the construction of pushouts in  $\text{mod}(\Sigma, \mathbf{E})$ ,  $D$  is the reflection of some  $P \in \text{mod}(\Sigma, \hat{\mathbf{C}})$  with



$$\begin{array}{ccc}
 A & \xrightarrow{f} & B \\
 \downarrow u & & \downarrow v \\
 C & \xrightarrow{g} & P
 \end{array}$$

is a pushout in  $\text{mod}(\Sigma, \hat{C})$ ,  $R\bar{g} = g$ , and  $f$  a monomorphism. Now, for each  $U \in \mathbf{C}$ ,

$$\begin{array}{ccc}
 AU & \xrightarrow{f_U} & BU \\
 \downarrow & & \downarrow \\
 CU & \xrightarrow{\bar{g}_U} & PU
 \end{array}$$

is a pushout in  $\text{mod } \Sigma$  with  $f_U$  a monomorphism. Then, by the hypothesis on  $\text{mod } \Sigma$ ,  $\bar{g}_U$  (for all  $U$ ), and hence  $\bar{g}$  are monomorphisms. Since  $R\bar{g} = g$  and  $R$  preserves monomorphisms,  $g$  is a monomorphism.

**4.3. Lemma.** *The category  $\text{mod}(\Sigma, \mathbf{E})$  has enough injectives iff, it is residually small and pushouts transfer monomorphisms.*

**Proof.** ( $\Rightarrow$ ) To show that  $\text{mod}(\Sigma, \mathbf{E})$  is residually small is to show it is essentially bounded, by Corollary 3.9. All the essential extensions of  $A \in \text{mod}(\Sigma, \mathbf{E})$  can be embedded in any injective extension of  $A$ , hence there exists a set of essential extensions for  $A$ . To prove the second part, let  $A \xrightarrow{f} B$  be a monomorphism and  $A \xrightarrow{g} C$  any homomorphism. Embedding  $C$  into an injective  $E$  we get the following commutative diagram

$$\begin{array}{ccc}
 A & \xrightarrow{f} & B \\
 \downarrow g & & \downarrow h \\
 C & \xrightarrow{j} & E
 \end{array}$$

where  $h$  is obtained because  $f$  is a monomorphism and  $E$  is injective. This, by an earlier remark, proves that pushouts transfer monomorphisms.

( $\Leftarrow$ ) For  $A \in \text{mod}(\Sigma, \mathbf{E})$ , take a maximal extension  $f: A \rightarrow E$  of  $A$ , which exists by Lemma 3.2 (ii). We claim that  $E$  is injective. To prove this, let  $g: B \rightarrow C$  be any monomorphism and  $h: B \rightarrow E$  any homomorphism. Form the following pushout:

$$\begin{array}{ccc}
 B & \xrightarrow{g} & C \\
 \downarrow h & & \downarrow v \\
 E & \xrightarrow{u} & P
 \end{array}$$

By hypothesis  $u$  is a monomorphism, and hence retractable by Corollary 3.4. This proves that  $E$  is injective.

**4.4. Proposition.** *The category  $\text{mod}(\Sigma, E)$  has enough injectives iff  $\text{mod } \Sigma$  has enough injectives.*

**Proof.** This follows from Propositions 3.10 and 4.2 and Lemma 4.3.

This result substantially improves a similar result by Howlett [5]. Here, we deal with quasi-equational classes of algebras rather than equational classes as [5] does, but more importantly, our proof does not use the points of the topos whereas [5] only proves this result for a Grothendieck topos with enough points. Moreover, Proposition 3.10 provides a positive answer to Howlett's question [5, p. 108] whether essential boundedness of  $\text{mod } \Sigma$  directly implies that of  $\text{mod}(\Sigma, E)$ .

## 5. Behaviour of injectivity in $\text{mod}(\Sigma, E)$

**5.1.** Banaschewski in [1] calls the notion of injectivity in a category  $\mathbf{K}$  *properly behaved* iff the following three propositions hold, which describe the relationship between essential boundedness, residual smallness and the existence of injective hulls in  $\mathbf{K}$ . Actually [1] deals with injectivity with respect to a more general type of morphism, but of course, here we only consider injectivity with respect to all monomorphisms.

(I) For any  $A \in \mathbf{K}$  the following conditions are equivalent:

- (I1)  $A$  is injective.
- (I2)  $A$  is an absolute retract.
- (I3)  $A$  has no proper essential extensions.

(E) Every  $A \in \mathbf{K}$  has an injective hull, unique up to isomorphisms.

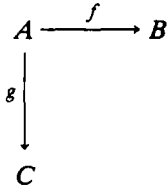
(H) For any monomorphism  $f: A \rightarrow B$ , the following conditions are equivalent:

- (H1)  $f: A \rightarrow B$  is an injective hull of  $A$ .
- (H2)  $f: A \rightarrow B$  is a maximal essential extension.
- (H3)  $f: A \rightarrow B$  is a minimal injective extension.

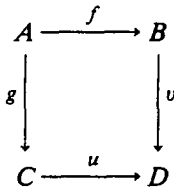
**5.2.** [1] also gives sufficient conditions for the proper behaviour of injectivity in  $\mathbf{K}$  as follows:

(E3) For any monomorphism  $f: A \rightarrow B$ , there exists a homomorphism  $g: B \rightarrow C$  with  $gf$  an essential monomorphism.

(E4) Any diagram



with  $f$  a monomorphism can be completed to a commutative diagram



such that  $u$  is a monomorphism.

(E5) Any direct limit of monomorphisms is a monomorphism.

(E6) The category  $\mathbf{K}$  is essentially bounded.

For  $\mathbf{K} = \text{mod}(\Sigma, \mathbf{E})$ , we now have the following counterpart of Proposition 5 in [1] for equational classes of algebras in  $\text{Ens}$ .

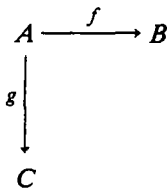
**5.3. Proposition.** *For  $\text{mod}(\Sigma, \mathbf{E})$ , the following are equivalent:*

- (i) *Injectivity is properly behaved.*
- (ii)  *$\text{Mod}(\Sigma, \mathbf{E})$  has enough injectives.*
- (iii)  *$\text{Mod}(\Sigma, \mathbf{E})$  is residually small and pushouts transfer monomorphisms.*
- (iv) *(E4) and (E6) are satisfied.*

**Proof.** (i)  $\Rightarrow$  (ii): By (E) in the definition of proper behaviour of injectivity.

(ii)  $\Rightarrow$  (iii): By Lemma 4.3.

(iii)  $\Rightarrow$  (iv): (E4) is trivial, by completing any diagram



to a pushout, and (E6) holds by Corollary 3.9.

(iv)  $\rightarrow$  (i): It remains to show (E3) and (E5). But Lemma 3.3 proves (E3), and (E5) is discussed in the proof of Lemma 3.2 (ii).

In particular, one has, by Proposition 4.4:

*Injectivity is properly behaved in  $\text{mod}(\Sigma, \mathbf{E})$  iff it is properly behaved in  $\text{mod } \Sigma$ .*

We conclude with a couple of comments on injectivity in  $\text{mod}(\Sigma, \mathbf{C})$ .  
 Since each of the functors

$$\text{mod}(\Sigma, \mathbf{E}) \xrightarrow{(V, -)} \text{mod } \Sigma \quad (V \in \mathbf{C})$$

has a left adjoint preserving monomorphisms, and by the well known fact that such a functor preserves injectives, if  $A \in \text{mod}(\Sigma, \mathbf{E})$  is injective, then so is each  $AV$  ( $V \in \mathbf{C}$ ). However, the converse of this is not true; for counter examples, in the case of abelian groups, the reader is referred to [2].

For certain  $\Sigma$ , one has characterizations of the injective  $A \in \text{mod } \Sigma$  by properties of  $A$  in terms of its elements of its elements and subsets, for example: divisibility for abelian groups, completeness for Boolean algebras, and completeness and Booleanness for distributive lattices. An obvious question to ask is to what extent, that is for what  $\mathbf{E}$ , such characterizations remain valid in  $\text{mod}(\Sigma, \mathbf{E})$ . The only case where anything is known about this is that of abelian groups: divisibility = injectivity for abelian groups in the category  $\text{Sh } \mathbf{L}$  of sheaves on a locale  $\mathbf{L}$  iff  $\mathbf{L}$  is Boolean [3].

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